

Quantum expanders – Random constructions & Applications

Based on joint works with:

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- 1 Introduction: classical and quantum expanders
- 2 Random constructions of expanders
- 3 Implications for random matrix product states

Classical expanders

G a d -regular graph on n vertices (d edges at each vertex).

A its (normalized) adjacency matrix, i.e. the $n \times n$ matrix s.t. $A_{kl} = e(k, l)/d$ for all $1 \leq k, l \leq n$.
number of edges between vertices k and l ↴

$\lambda_1(A), \dots, \lambda_n(A)$ eigenvalues of A , ordered s.t. $|\lambda_1(A)| \geq \dots \geq |\lambda_n(A)|$.

G regular $\implies \lambda_1(A) = 1$ with associated eigenvector the uniform probability $u = (1/n, \dots, 1/n)$.

The *spectral expansion parameter* of G is $\lambda(G) := |\lambda_2(A)|$.

Observation: $\lambda(G) = |\lambda_1(A - J)|$, where J is the adjacency matrix of the *complete graph* on n vertices, i.e. the matrix whose entries are all equal to $1/n$.

$\longrightarrow \lambda(G)$ is a distance measure between G and the complete graph.

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 $\rightarrow \lambda(G)$ is a distance measure between G and the complete graph.

Definition [Classical expander]

A d -regular graph G on n vertices is an *expander* if it is sparse (i.e. $d \ll n$) and spectrally expanding (i.e. $\lambda(G) \ll 1$).

$\rightarrow G$ is both 'economical' and 'resembling' the complete graph.

For instance, a random walk supported on G converges fast to equilibrium.

Indeed, for any probability p on $\{1, \dots, n\}$, $\forall q \in \mathbf{N}$, $\|A^q p - u\|_1 \leq \sqrt{n} \|A^q p - u\|_2 \leq \sqrt{n} \lambda(G)^q$.
exponential convergence, at rate $|\log \lambda(G)|$ \leftarrow

Quantum analogue of the transition matrix associated to a regular graph

Classical - Quantum correspondence:

- $p \in \mathbf{R}^n$ probability vector $\longleftrightarrow \rho \in \mathcal{M}_n(\mathbf{C})$ density operator (PSD and trace 1 operator).
- $A : \mathbf{R}^n \rightarrow \mathbf{R}^n$ transition matrix $\longleftrightarrow \Phi : \mathcal{M}_n(\mathbf{C}) \rightarrow \mathcal{M}_n(\mathbf{C})$ quantum channel (CPTP map).
- G regular: A leaves $\mathbf{1}$ invariant $\longleftrightarrow \Phi$ unital: Φ leaves I/n invariant.

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Question: What is the analogue of the degree?

Given a CP map Φ on $\mathcal{M}_n(\mathbf{C})$, its *Kraus representation* is:

$$\Phi: X \in \mathcal{M}_n(\mathbf{C}) \mapsto \sum_{i=1}^d K_i X K_i^* \in \mathcal{M}_n(\mathbf{C}), \text{ where } K_1, \dots, K_d \in \mathcal{M}_n(\mathbf{C}). \quad (*)$$

\downarrow Kraus operators of Φ

The minimal d s.t. Φ can be written as $(*)$ is the *Kraus rank* of Φ (it is always at most n^2).

[Note: Φ is TP iff $\sum_{i=1}^d K_i^* K_i = I$. Φ is unital iff $\sum_{i=1}^d K_i K_i^* = I$.]

Answer: The analogue of the degree is the Kraus rank.

- G a d -regular graph: If $|\text{supp}(p)| = 1$, then $|\text{supp}(Ap)| \leq d$.
- Φ a Kraus rank d unital quantum channel: If $\text{rk}(\rho) = 1$, then $\text{rk}(\Phi(\rho)) \leq d$.

\longrightarrow Both quantify the 1-iteration spreading.

Φ a Kraus rank d unital quantum channel on $\mathcal{M}_n(\mathbf{C})$.

$\lambda_1(\Phi), \dots, \lambda_{n^2}(\Phi)$ eigenvalues of Φ , ordered s.t. $|\lambda_1(\Phi)| \geq \dots \geq |\lambda_{n^2}(\Phi)|$.

Φ unital $\implies \lambda_1(\Phi) = 1$ with associated eigenstate the maximally mixed state I/n .

The *spectral expansion parameter* of Φ is $\lambda(\Phi) := |\lambda_2(\Phi)|$.

Observation: $\lambda(\Phi) = |\lambda_1(\Phi - \Pi)|$, where Π is the *maximally mixing channel* on $\mathcal{M}_n(\mathbf{C})$, i.e.

$\Pi : X \in \mathcal{M}_n(\mathbf{C}) \mapsto \text{Tr}(X) I/n \in \mathcal{M}_n(\mathbf{C})$.

$\implies \lambda(\Phi)$ is a distance measure between Φ and the maximally mixing channel.

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Definition [Quantum expander]

A Kraus rank d unital quantum channel Φ on $\mathcal{M}_n(\mathbf{C})$ is an *expander* if it is little noisy (i.e. $d \ll n^2$) and spectrally expanding (i.e. $\lambda(\Phi) \ll 1$).

$\implies \Phi$ is both 'economical' and 'resembling' the maximally mixing channel.

For instance, the dynamics associated to Φ converges fast to equilibrium.

Indeed, for any state ρ on \mathbf{C}^n , $\forall q \in \mathbf{N}$, $\|\Phi^q(\rho) - I/n\|_1 \leq \sqrt{n} \|\Phi^q(\rho) - I/n\|_2 \leq \sqrt{n} \lambda(\Phi)^q$.
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2 Random constructions of expanders

- Classical case
- Quantum case

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- Definitions and motivations
- Decay of correlations in random translation-invariant matrix product states

Constructions of optimal classical expanders

Fact: For any d -regular graph G on n vertices, $\lambda(G) \geq 2\sqrt{d-1}/d - o_n(1)$.

→ G is called a *Ramanujan graph* if it is an optimal expander, i.e. $\lambda(G) \leq 2\sqrt{d-1}/d$.

Question: Do Ramanujan graphs exist?

- 1 Explicit constructions of exactly Ramanujan graphs only for $d = p^m + 1$, p prime.
- 2 Random constructions of almost Ramanujan graphs for all d .
- 3 Existence of exactly Ramanujan graphs for all d .

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In fact, for large n , almost all regular graphs are almost Ramanujan:

Theorem [Uniform random regular graph (Friedman, Bordenave)]

Fix $d \in \mathbf{N}$. Let G be uniformly distributed on the set of d -regular graphs on n vertices.

Then, for all $\varepsilon > 0$, $\mathbf{P}\left(\lambda(G) \leq \frac{2\sqrt{d-1}}{d} + \varepsilon\right) = 1 - o_n(1)$.

Remarks: ↗ permutation model

- First proven for a simpler model of random regular graphs: for d even, pick $\sigma_1, \dots, \sigma_{d/2} \in \mathcal{S}_n$ independent uniformly distributed and let G have edges $\{(k, \sigma_i(k)), (k, \sigma_i^{-1}(k))\}_{1 \leq k \leq n, 1 \leq i \leq d/2}$.
- Remains true for d_n growing with n , up to a constant multiplicative factor (Cook/Goldstein/Johnson, Tikhomirov/Youssef): $\mathbf{P}(\lambda(G) \leq C/\sqrt{d_n} + \varepsilon) = 1 - o_n(1)$.

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Constructions of optimal quantum expanders

Fact: For any Kraus rank d unital quantum channel Φ on $\mathcal{M}_n(\mathbf{C})$, $\lambda(\Phi) \geq c/\sqrt{d}$.

→ Φ is an optimal expander if $\lambda(\Phi) \leq C/\sqrt{d}$.

Question: Do optimal quantum expanders exist?

First attempts at exhibiting explicit constructions (inspired by classical ones): not optimal.

→ What about random constructions?

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Question: How to sample a unital quantum channel randomly?

Idea: Pick random Kraus operators $K_1, \dots, K_d \in \mathcal{M}_n(\mathbf{C})$, under the constraint
$$\begin{cases} \sum_{i=1}^d K_i^* K_i = I \\ \sum_{i=1}^d K_i K_i^* = I \end{cases} .$$

Let $\Phi : X \in \mathcal{M}_n(\mathbf{C}) \mapsto \sum_{i=1}^d K_i X K_i^* \in \mathcal{M}_n(\mathbf{C})$ be the associated random unital quantum channel.

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Theorem [Independent paired Haar unitaries as Kraus operators (Hastings, Pisier)]

Fix $d \in \mathbf{N}$ even. Pick $U_1, \dots, U_{d/2} \in \mathcal{M}_n(\mathbf{C})$ independent Haar unitaries. Let $K_i = U_i/\sqrt{d}$, $1 \leq i \leq d/2$. The random CP map Φ associated to the K_i 's, K_i^* 's is TP and unital by construction.

Then, for all $\varepsilon > 0$, $\mathbf{P} \left(\lambda(\Phi) \leq \frac{2\sqrt{d-1}}{d} + \varepsilon \right) = 1 - o_n(1)$.

Remarks:

- Optimal for quantum channels Φ with d unitary Kraus operators, for which $\lambda(\Phi) \geq 2\sqrt{d-1}/d$.
- Same result, up to a constant multiplicative factor, for d independent unitary Kraus operators.

More random examples of optimal quantum expanders

Question: Can the previous result be extended to other random models? And to a regime where d is not fixed but grows with n ?

Difficulty: Imposing that Φ is both TP and unital is very constraining.

However, the definition of expander can be extended to 'close to unital' quantum channels, whose fixed point ρ_* has a large entropy: $S(\rho_*) \geq \alpha S(I/n) = \alpha \log n$, for some $0 < \alpha < 1$.

[Note: We now have $\lambda(\Phi) = |\lambda_1(\Phi - \Pi_{\rho_*})|$, where $\Pi_{\rho_*} : X \in \mathcal{M}_n(\mathbf{C}) \mapsto \text{Tr}(X)\rho_* \in \mathcal{M}_n(\mathbf{C})$.]

Classical analogy: Relaxation of the exact regularity condition, e.g. to look at Erdős-Rényi graphs.

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Theorem [Independent Gaussians as Kraus operators (Lancien/Pérez-García)]

Pick $G_1, \dots, G_d \in \mathcal{M}_n(\mathbf{C})$ independent Gaussian matrices. Let $\tilde{K}_i = G_i/\sqrt{d}$, $1 \leq i \leq d$.
↳ i.i.d. Gaussian entries (mean 0 and variance $1/n$)

The random CP map $\tilde{\Phi}$ associated to the \tilde{K}_i 's is not TP but almost: $\mathbf{P}(\Sigma := \sum_{i=1}^d \tilde{K}_i^* \tilde{K}_i \simeq I) \simeq 1$.

With $K_i = \tilde{K}_i \Sigma^{-1/2}$, $1 \leq i \leq d$, the random CP map Φ associated to the K_i 's is TP by construction.

Then, $\mathbf{P}\left(S(\rho_*) \geq \log n - \frac{C'}{\sqrt{d}} \text{ and } \lambda(\Phi) \leq \frac{C}{\sqrt{d}}\right) \geq 1 - e^{-cn}$, for $C, C', c > 0$ constants.

Remark: Other model that was proven to be a.s. an optimal expander as n grows (for d fixed): blocks of a Haar isometry $V : \mathbf{C}^n \hookrightarrow \mathbf{C}^n \otimes \mathbf{C}^d$ as Kraus operators (González-Guillén/Junge/Nechita).

How much can the previous examples be generalized?

Theorem [Independent general random matrices as Kraus operators (Lancien/Youssef)]

- 1 Let $A \in \mathcal{M}_n(\mathbf{R})$ be a doubly stochastic matrix s.t. $|\lambda_2(A)| \leq \frac{C}{\sqrt{d}}$, with $d \geq (\log n)^4$.
E.g. A the adjacency matrix of a d -regular graph G on n vertices s.t. $\lambda(G) \leq \frac{C}{\sqrt{d}}$.
- 2 Let $W \in \mathcal{M}_n(\mathbf{C})$ be a random matrix with independent centered entries, s.t.
 $\forall 1 \leq k, l \leq n, \mathbf{E}|W_{kl}|^2 = A_{kl}$ and $(\mathbf{E}|W_{kl}|^{2p})^{1/p} \leq C' p^\beta A_{kl}, p \in \mathbf{N}$.
[$\beta = 0$: bounded entries. $\beta = 1$: sub-Gaussian entries. $\beta = 2$: sub-exponential entries.]
- 3 Pick $W_1, \dots, W_d \in \mathcal{M}_n(\mathbf{C})$ independent copies of W . Let $K_i = \frac{W_i}{\sqrt{d}}, 1 \leq i \leq d$, and Φ be the random CP map with the K_i 's as Kraus operators.

Then, Φ is on average TP and unital, and s.t. $\mathbf{E}\lambda(\Phi) \leq \frac{C''}{\sqrt{d}}$.

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Then, Φ is on average TP and unital, and s.t. $\mathbf{E}\lambda(\Phi) \leq \frac{C''}{\sqrt{d}}$.

Interest: Constructing a random optimal quantum expander from any optimal classical expander.

→ Optimal quantum expanders can be obtained from random Kraus operators which are sparse and whose entries have any distribution following the moments' growth assumption.

Proof idea to show that $\mathbf{E}\lambda(\Phi) \leq C/\sqrt{d}$

Goal: In all cases, we want to upper bound $\mathbf{E}|\lambda_2(\Phi)| = \mathbf{E}|\lambda_1(\Phi - \Pi_{\rho^*})|$.

First step: Upper bound $\mathbf{E}|\lambda_1(\Phi - \mathbf{E}(\Phi))|$ (and then show that $\mathbf{E}(\Phi)$ is close to Π_{ρ^*}).

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• Observation 1: $|\lambda_1(\Psi)| \leq s_1(\Psi) = \|\Psi\|_\infty$.

• Observation 2: $\|\Psi\|_\infty = \|M_\Psi\|_\infty$, where for $\Psi : X \mapsto \sum_{i=1}^d K_i X L_i^*$, $M_\Psi = \sum_{i=1}^d K_i \otimes \bar{L}_i$.

[Identification $\Psi : \mathcal{M}_n(\mathbf{C}) \rightarrow \mathcal{M}_n(\mathbf{C}) \equiv M_\Psi : \mathbf{C}^n \otimes \mathbf{C}^n \rightarrow \mathbf{C}^n \otimes \mathbf{C}^n$ preserves the operator norm.]

→ We want to upper bound $\mathbf{E}\| \underbrace{M_\Phi - \mathbf{E}(M_\Phi)}_{=: X} \|_\infty$, where $M_\Phi = \sum_{i=1}^d K_i \otimes \bar{K}_i$ with the K_i 's random.

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↳ Haar unitaries, Gaussians, blocks of Haar isometry

• For concrete models, this can be done by a moments' method:

By Jensen's inequality, we have: $\forall p \in \mathbf{N}$, $\mathbf{E}\|X\|_\infty \leq \mathbf{E}\|X\|_p \leq (\mathbf{E}\text{Tr}|X|^p)^{1/p}$.

The term on the r.h.s. can be estimated and provides a good upper bound for $p \simeq n^\gamma$.

↳ by Weingarten or Wick calculus

• For the general case, we use recent results on estimating the operator norm of random matrices with dependencies and non-homogeneity (Bandeira/Boedihardjo/van Handel, Brailovskaya/van Handel):

Setting $X = \sum_{i=1}^d Z_i$, with $Z_i := K_i \otimes \bar{K}_i - \mathbf{E}(K_i \otimes \bar{K}_i)$, $1 \leq i \leq d$, we have for $p \simeq \log n$,

$$\mathbf{E}\|X\|_\infty \lesssim \|\mathbf{E}(X X^*)\|_\infty^{1/2} + \|\mathbf{E}(X^* X)\|_\infty^{1/2} + (\log n)^{3/2} \|\mathbf{Cov}(X)\|_\infty^{1/2} + (\log n)^2 \left(\sum_{i=1}^d \mathbf{E}\text{Tr}|Z_i|^p \right)^{1/p}.$$

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- **Definitions and motivations**
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Small subset of 'physically relevant' states of many-body quantum systems

Curse of dimensionality: Exponential growth of system's dimension with number of subsystems. However, 'physically relevant' states of many-body quantum systems are often well approximated by so-called *tensor network states (TNS)*, which form a small subset of the global state space.

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Example: A *matrix product state (MPS)* on $(\mathbf{C}^d)^{\otimes M}$ is a pure state $\chi \in (\mathbf{C}^d)^{\otimes M}$ of the form

$$\chi = \sum_{i_1, \dots, i_M=1}^d \text{Tr} \left(K_{i_1}^{(1)} \dots K_{i_M}^{(M)} \right) |e_{i_1}\rangle \otimes \dots \otimes |e_{i_M}\rangle, \text{ where } K_1^{(\ell)}, \dots, K_d^{(\ell)} \in \mathcal{M}_n(\mathbf{C}), 1 \leq \ell \leq M.$$

→ Such state is described by Mdn^2 parameters, which is linear rather than exponential in M .

[Vocabulary: d is the *physical dimension*. n is the *bond dimension*.]

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Fact: On a 1D system (M subsystems disposed on a line), the *ground state of a gapped local Hamiltonian* is well approximated by an MPS (Hastings, Landau/Vazirani/Vidick...)

spectral gap lower bounded by a constant independent of M ←

composed of terms which act non-trivially only on nearby sites ←

→ In condensed-matter physics, MPS are used as Ansatz in ground energy computations: optimization over a manageable number of parameters, even for large M .

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Model of random translation-invariant MPS

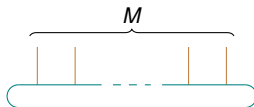
Idea: Pick $K_1, \dots, K_d \in \mathcal{M}_n(\mathbf{C})$ at random.

Let $\chi \in (\mathbf{C}^d)^{\otimes M}$ be the corresponding *random translation-invariant MPS*, i.e.

$$\chi = \sum_{i_1, \dots, i_M=1}^d \text{Tr}(K_{i_1} \cdots K_{i_M}) |e_{i_1}\rangle \otimes \cdots \otimes |e_{i_M}\rangle.$$



$$K = \sum_{i=1}^d K_i \otimes |e_i\rangle \in \mathcal{M}_n(\mathbf{C}) \otimes \mathbf{C}^d$$



$$\chi = \sum_{i_1, \dots, i_M=1}^d \text{Tr}(K_{i_1} \cdots K_{i_M}) |e_{i_1}\rangle \otimes \cdots \otimes |e_{i_M}\rangle \in (\mathbf{C}^d)^{\otimes M}$$

Model of random translation-invariant MPS

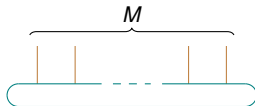
Idea: Pick $K_1, \dots, K_d \in \mathcal{M}_n(\mathbf{C})$ at random.

Let $\chi \in (\mathbf{C}^d)^{\otimes M}$ be the corresponding *random translation-invariant MPS*, i.e.

$$\chi = \sum_{i_1, \dots, i_M=1}^d \text{Tr}(K_{i_1} \cdots K_{i_M}) |e_{i_1}\rangle \otimes \cdots \otimes |e_{i_M}\rangle.$$



$$K = \sum_{i=1}^d K_i \otimes |e_i\rangle \in \mathcal{M}_n(\mathbf{C}) \otimes \mathbf{C}^d$$



$$\chi = \sum_{i_1, \dots, i_M=1}^d \text{Tr}(K_{i_1} \cdots K_{i_M}) |e_{i_1}\rangle \otimes \cdots \otimes |e_{i_M}\rangle \in (\mathbf{C}^d)^{\otimes M}$$

Associated *transfer operator*. $T = \sum_{i=1}^d K_i \otimes \bar{K}_i \in \mathcal{M}_n(\mathbf{C}) \otimes \mathcal{M}_n(\mathbf{C})$,

obtained by contracting the d -dimensional indices of K and \bar{K} .

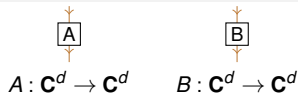


Observation: $T = \sum_{i=1}^d K_i \otimes \bar{K}_i$ is the matrix version of the CP map $\Phi_T : X \mapsto \sum_{i=1}^d K_i X K_i^*$.
[Identification $T : \mathbf{C}^n \otimes \mathbf{C}^n \rightarrow \mathbf{C}^n \otimes \mathbf{C}^n \equiv \Phi_T : \mathcal{M}_n(\mathbf{C}) \rightarrow \mathcal{M}_n(\mathbf{C})$.]

In particular, both have the same spectrum.

Correlations in an MPS

Let A, B be 1-site observables, i.e. observables on \mathbf{C}^d .



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$$\begin{array}{c} \downarrow \\ \boxed{A} \\ \downarrow \\ A : \mathbf{C}^d \rightarrow \mathbf{C}^d \end{array} \qquad \begin{array}{c} \downarrow \\ \boxed{B} \\ \downarrow \\ B : \mathbf{C}^d \rightarrow \mathbf{C}^d \end{array}$$

Compute the value on the MPS χ of the observable $A_1 \otimes I_q \otimes B_1 \otimes I_{M-q-2}$, i.e.

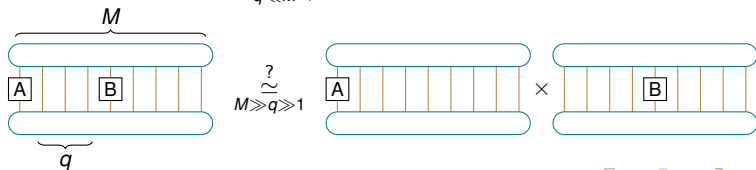
$$v_\chi(A, B, q) := \frac{\langle \chi | A_1 \otimes I_q \otimes B_1 \otimes I_{M-q-2} | \chi \rangle}{\langle \chi | \chi \rangle}.$$

Compare it to the product of the values on χ of $A_1 \otimes I_{M-1}$ and $I_{q+1} \otimes B_1 \otimes I_{M-q-2}$, i.e.

$$v_\chi(A)v_\chi(B) := \frac{\langle \chi | A_1 \otimes I_{M-1} | \chi \rangle \langle \chi | I_{q+1} \otimes B_1 \otimes I_{M-q-2} | \chi \rangle}{\langle \chi | \chi \rangle^2}.$$

Correlations in the MPS χ : $\gamma_\chi(A, B, q) := |v_\chi(A, B, q) - v_\chi(A)v_\chi(B)|$.

Question: Do we have $\gamma_\chi(A, B, q) \xrightarrow{q \ll M \rightarrow \infty} 0$? And if so, at which speed?



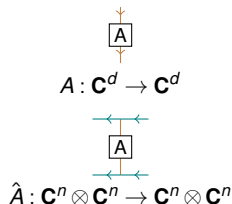
Correlation length in an MPS and spectrum of its transfer operator

$\chi \in (\mathbf{C}^d)^{\otimes M}$ an MPS.

T its associated transfer operator on $\mathbf{C}^n \otimes \mathbf{C}^n$.

Its correlation function can be re-written as:

$$\gamma_\chi(A, B, q) = \left| \frac{\text{Tr}(\hat{A}T^q\hat{B}T^{M-q-2})}{\text{Tr}(T^M)} - \frac{\text{Tr}(\hat{A}T^{M-1})\text{Tr}(\hat{B}T^{M-1})}{(\text{Tr}(T^M))^2} \right|.$$



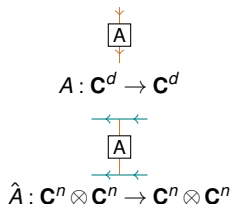
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Important consequence:

$\lambda_1(T), \lambda_2(T), \dots$ eigenvalues of T (with multiplicities), ordered s.t. $|\lambda_1(T)| > |\lambda_2(T)| > \dots$.

Set $\Delta(T) := |\lambda_1(T)| - |\lambda_2(T)|$ and $\varepsilon(T) = |\lambda_2(T)|/|\lambda_1(T)|$.

$$\gamma_\chi(A, B, q) \leq C \left(\frac{\text{Tr}(T)}{\Delta(T)} \right)^2 \varepsilon(T)^q \|A\|_\infty \|B\|_\infty.$$

→ Correlations between two 1-site observables decay exponentially with the distance separating the two sites, at a rate $\tau(\chi) = |\log \varepsilon(T)|$.

Correlation length in the MPS χ : $\xi(\chi) := 1/\tau(\chi) = 1/|\log \varepsilon(T)|$.

Conclusion: Estimating $\xi(\chi)$ boils down to estimating $|\lambda_1(\Phi_T)|$ and $|\lambda_2(\Phi_T)|$.

Decay of correlations in random translation-invariant MPS

Examples of distribution for $K_1, \dots, K_d \in \mathcal{M}_n(\mathbf{C})$:

- 1 $K_i = W_i/\sqrt{d}$, $1 \leq i \leq d$, where the W_i 's are i.i.d. matrices, with independent centered entries having variance profile a doubly stochastic matrix A s.t. $|\lambda_2(A)| \leq C/\sqrt{d}$.
E.g. Gaussians with mean 0 and variance $1/n$.
- 2 $K_i = U_i/\sqrt{d}$, $1 \leq i \leq d$, where the U_i 's are i.i.d. Haar unitaries.
- 3 $K_i = V_i$, $1 \leq i \leq d$, where $V = \sum_{i=1}^d V_i \otimes |e_i\rangle$ is a Haar isometry.

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Theorem [Correlation length of a random translation-invariant MPS (Lancien/Pérez-García)]

Let $\chi \in (\mathbf{C}^d)^{\otimes M}$ be a random translation-invariant MPS, with associated $K_1, \dots, K_d \in \mathcal{M}_n(\mathbf{C})$ sampled according to one of the models above.

For large n , its correlation length is typically upper bounded by $2/\log d$.

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For large n , its correlation length is typically upper bounded by $2/\log d$.

Question: What about more complicated models, where the random MPS has a local symmetry?

For instance: Let G be a compact group with unitary representation $U : G \rightarrow \mathcal{M}_d(\mathbf{C})$.

Assume that $\chi \in (\mathbf{C}^d)^{\otimes M}$ is s.t. $U_g^{\otimes M} \chi = \chi$ for all $g \in G$.

This means that $K_1, \dots, K_d \in \mathcal{M}_n(\mathbf{C})$ are s.t. $P_g K_i P_g^* = K_i$ for all $g \in G$, where $P : G \rightarrow \mathcal{M}_n(\mathbf{C})$ is a projective representation (Schuch/Pérez-García/Cirac).

$$\hookrightarrow P_g P_h = e^{i\omega(g,h)} P_{gh}$$

\rightarrow Given $K_1, \dots, K_d \in \mathcal{M}_n(\mathbf{C})$ random, set $\hat{K}_i := \Phi_G(K_i)$, where $\Phi_G(X) = \mathbf{E}_{g \in G}[P_g X P_g^*]$.

One-slide summary of “Implications for random matrix product states”

Matrix product states (MPS) form a subset of many-body quantum states.

They are particularly useful because:

- They admit an *efficient description* (number of parameters that scales linearly rather than exponentially with the number of subsystems).
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with the distance separating the sites ←
between observables measured on distinct sites ←

Main result: Random MPS typically have correlations that decay exponentially fast, with a *small correlation length*.

Proof strategy: Observe that the correlation length is given by $1/|\log \lambda(\Phi)|$ for Φ a random quantum channel associated to the random MPS (its so-called *transfer operator*).

Some perspectives

- What about *explicit constructions* of optimal quantum expanders? Would be important in practice (quantum cryptography, error correction, condensed matter physics, etc.)

Previously known constructions required a large amount of randomness. First step towards *derandomization*: taking Kraus operators as sparse Rademacher matrices works as well.

Other direction to explore: unitary Kraus operators sampled according to a 'simple' measure that still 'resembles' the uniform one. E.g. an approximate t -design measure or the uniform measure on a subgroup of the unitary group (Bordenave/Collins).

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- What about identifying the *full spectral distribution* of random quantum channels?

Known: For a random quantum channel $\Phi : \mathcal{M}_n(\mathbf{C}) \rightarrow \mathcal{M}_n(\mathbf{C})$ sampled according to various models, the eigenvalues of $\Phi - \Pi_{\rho_*}$ are typically inside a disc of radius C/\sqrt{d} for large n .
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- Do the results about the typical spectral gap of random quantum channels remain true when we impose *extra symmetries* on the model?
- What about looking at other, related, notions of expansions, such as *geometric* ones (Bannink/Briët/Labib/Maassen) or *linear-algebraic* ones (Li/Qiao/Wigderson/Wigderson/Zhang)?

- **T. Bannink, J. Briët, F. Labib, H. Maassen.** Quasirandom quantum channels. 2020.
- **J. Friedman.** A proof of Alon's second eigenvalue conjecture. 2008.
- **A. Bandeira, M. Boedihardjo, R. van Handel.** Matrix concentration inequalities and free probability. 2021.
- **C. Bordenave.** A new proof of Friedman's second eigenvalue theorem and its extension to random lifts. 2020.
- **C. Bordenave, B. Collins.** Strong asymptotic freeness for independent uniform variables on compact groups associated to non-trivial representations. 2020.
- **T. Brailovskaya, R. van Handel.** Universality and sharp matrix concentration inequalities. 2022.
- **N. Cook, L. Goldstein, T. Johnson.** Size biased couplings and the spectral gap for random regular graphs. 2018.
- **C. González-Guillén, M. Junge, I. Nechita.** On the spectral gap of random quantum channels. 2018.
- **M.B. Hastings.** Solving gapped Hamiltonians locally. 2006.
- **M.B. Hastings.** Random unitaries give quantum expanders. 2007.
- **C. Lancien, D. Pérez-García.** Correlation length in random MPS and PEPS. 2021.
- **C. Lancien, P. Youssef.** A note on quantum expanders. 2023.
- **Z. Landau, U. Vazirani, T. Vidick.** A polynomial-time algorithm for the ground state of 1D gapped local Hamiltonians. 2015.
- **Y. Li, Y. Qiao, A. Wigderson, Y. Wigderson, C. Zhang.** On linear-algebraic notions of expansion. 2022.
- **G. Pisier.** Quantum expanders and geometry of operator spaces. 2014.
- **N. Schuch, D. Pérez-García, J.I. Cirac.** Classifying quantum phases using MPS and PEPS. 2011.
- **K. Tikhomirov, P. Youssef.** The spectral gap of dense random regular graphs. 2019.